

Das Auswahlaxiom A.8

Sprache: Englisch

Definition: A Boolean algebra is an algebraic structure $(B, +, \cdot, -, 0, 1)$ where B is a non-empty set, "+" and " $\cdot$ " are two binary operations (called Boolean sum and product), " $-$ " is a unary operation (called complement), and $0, 1$ are two constants.

For all $u, v, w \in B$, the Boolean operations satisfy the following axioms:

1) Commutativity

$$u + v = v + u$$

$$u \cdot v = v \cdot u$$

2) Associativity

$$u + (v + w) = (u + v) + w$$

$$u \cdot (v \cdot w) = (u \cdot v) \cdot w$$

3) Distributivity

$$u \cdot (v + w) = (u \cdot v) + (u \cdot w)$$

$$u + (v \cdot w) = (u + v) \cdot (u + w)$$

4) Absorption

$$u \cdot (u + v) = u$$

$$u + (u \cdot v) = u$$

5) Complementation

$$u + (-u) = 1$$

$$u \cdot (-u) = 0$$

Definition: An algebra of sets is a collection $\mathcal{S}$ of subsets of a given set S s.t.,

$$\bullet S \in \mathcal{S}$$

$$\bullet \forall X, Y \in \mathcal{S} : S \setminus (X \cap Y) \in \mathcal{S}$$

The last condition means that $\mathcal{S}$ is closed under unions, intersections and complements

Example:

An algebra of sets $\mathcal{S} \subseteq \mathcal{P}(S)$, for a given set S , is a Boolean algebra with:

Boolean sum: $\cup$

Boolean product: $\cap$

Complement: $S \setminus \cdot$ (or $(\cdot)^c$)

$0 = \emptyset$, $1 = S$

In particular, for any set S ,

$(\mathcal{P}(S), \cup, \cap, (\cdot)^c, \emptyset, S)$ is a Boolean algebra

Exercises: From the axioms of a Boolean algebra, we can derive (in green, the meaning for our example)

i) $0 + 0 = 0 \cdot 0 = -(-0) = 0$

$$A \cup A = A \cap A = (A^c)^c = A$$

ii) $0 + 0 = 0$

$$A \cup \emptyset = A$$

iii) $0 \cdot 0 = 0$

$$A \cap \emptyset = \emptyset$$

iv) $0 + 1 = 1$

$$A \cup S = S$$

v) $0 \cdot 1 = 0$

$$A \cap S = A$$

Proposition:

1) The complement of an element $0 \in B$ is unique

2) $-0 = 1$ ($\emptyset^c = S$)

Proof:

1) Assume not, i.e. $\exists -0 \neq v \in B$ for $0 \in B$ st.

$$0 + (-0) = 1$$

$$0 \cdot (-0) = 0$$

$$0 + v = 1$$

$$0 \cdot v = 0$$

Then,

$$\begin{aligned}v &= v \cdot 1 \\&= v \cdot (u + (-u)) \\&\stackrel{3)+}{=} (v \cdot u) + (v \cdot (-u)) \\&= 0 + (v \cdot (-u)) \\&= (u \cdot (-u)) + (v \cdot (-u)) \\&\stackrel{3)+}{=} -u \cdot (u + v) \\&= -u \cdot 1 \\&= -u\end{aligned}$$

2) Using 1), it suffices to show:

• $0 + 1 = 1$, which is true by iv)

• $0 \cdot (-1) = 0$, which is true by iii)

Proposition: Furthermore, we can show the

De Morgan laws. Let $u, v \in B$, then.

1) $-(u + v) = (-u) \cdot (-v)$ $((A \cup B)^c = A^c \cap B^c)$

2) $-(u \cdot v) = (-u) + (-v)$ $((A \cap B)^c = A^c \cup B^c)$

Proof:

$$\begin{aligned}1) & (u + v) + ((-u) \cdot (-v)) \\&\stackrel{3)+}{=} ((u + v) + (-u)) \cdot ((u + v) + (-v)) \\&= (u + v + (-u)) \cdot (u + v + (-v)) \\&\stackrel{1)+5)}{=} (1 + v) \cdot (u + 1) \\&\stackrel{iv)}{=} 1 \cdot 1 \\&\stackrel{v)}{=} 1 \quad (1)\end{aligned}$$

$$\begin{aligned}& (u + v) \cdot ((-u) \cdot (-v)) \\&\stackrel{3)}{=} u \cdot ((-u) \cdot (-v)) + v \cdot ((-u) \cdot (-v)) \\&\stackrel{1)+2)}{=} (u \cdot (-u)) \cdot (-v) + (v \cdot (-v)) \cdot (-u) \\&\stackrel{5)}{=} 0 \cdot (-v) + 0 \cdot (-u) \\&\stackrel{ii)+iii)}{=} 0 \quad (2)\end{aligned}$$

By (1) and (2) and the fact that the complement is unique, we obtain:

$$-(\neg \neg \varphi) = (\neg \neg) \cdot (\neg \varphi)$$

2) Left to the reader

Definition. Let $\mathcal{L}$ be a first-order language and let S be the set of all $\mathcal{L}$ -sentences. We define an equivalence relation " $\sim$ " on S by:

$$\varphi \sim \psi \Leftrightarrow \vdash \varphi \leftrightarrow \psi, \text{ i.e.}$$

$\varphi \sim \psi$ iff, and only iff, φ is equivalent to ψ (in the logic sense)

The Lindebaum algebra is the set $B := S/\sim$ of all equivalence classes $[\varphi]$ with:

$$\begin{aligned} \cdot \quad [\varphi] + [\psi] &= [\varphi \vee \psi] & \cdot \quad 0 &= [\varphi \wedge \neg \varphi] \\ \cdot \quad [\varphi] \cdot [\psi] &= [\varphi \wedge \psi] & \cdot \quad 1 &= [\varphi \vee \neg \varphi] \\ \cdot \quad -[\varphi] &= [\neg \varphi] \end{aligned}$$

Define now

$$u \leq v \Leftrightarrow u \cdot (-v) = 0$$

The " $\leq$ " is then a partial ordering on B and we have:

$$u \leq v \Leftrightarrow u + v = v \Leftrightarrow u \vee v = v$$

$$\text{Note: } [\varphi] \leq [\psi] \Leftrightarrow \vdash \varphi \rightarrow \psi$$

Remarks:

• W.r.t " $\leq$ ", 1 is the greatest element of B , 0 is the least element

• For $u, v \in B$

$u + v$ is the least upper bound of $\{u, v\}$

$u \cdot v$ is the greatest lower bound of $\{u, v\}$

- We can define an additional operation " $\oplus$ " on B :

$$u \oplus v = (u \cdot (-v)) + (v \cdot (-u))$$

Then,

$$\begin{aligned} u \oplus u &= (u \cdot (-u)) + (u \cdot (-u)) \\ &\stackrel{D}{=} u \cdot (-u) \\ &\stackrel{S}{=} 0 \end{aligned}$$

So, every u is its own (and unique) inverse w.r.t. $\oplus$.

- $(B, \oplus, \cdot, 0, 1)$ is a ring (proof left to the reader)

Let $(B, +, \cdot, -, 0, 1)$ be a Boolean algebra.

Def: An ideal I in B is a non-empty proper subset of B with the following properties:

- $0 \in I$ but $1 \notin I$
- If $u \in I$ and $v \in I$, then $u+v \in I$
- For all $w \in B$ and all $u \in I$, $w \cdot u \in I$

Remark: The third condition is equivalent to $w \in B, u \in I, w \leq u \Rightarrow w \in I$

Proof: Since $w \leq u$, $w = w \cdot u \in I$ $\square$

We consider the Boolean algebra $(\mathcal{P}(w), \cup, \cap, -, \emptyset, w)$.

Claim: The set of all finite subsets of w , denoted by I_0 , is an ideal

Proof: $\emptyset \in I_0$, $w \notin I_0$

- Let u, v be finite subsets of w . Then $u \cup v$ is also finite
- Let $w \in \mathcal{P}(w)$, $u \in I_0$. Then $w \cap u \subseteq u$ and is thus finite $\square$

Def: We call this ideal the Fréchet ideal

Def: A filter F in B is a non-empty proper subset of B with the following properties:

- $0 \notin F$ but $1 \in F$
- If $u \in F$ and $v \in F$, then $u \cdot v \in F$
- For all $w \in B$ and all $u \in F$, $w+u \in F$

Remark: The last condition is equivalent to $w \in B, u \in F, w \geq u \Rightarrow w \in F$

Proof: Since $u \leq w$, $w = u+w \in F$ $\square$

Remark: Note that the notion of filter is the dual notion of the one of ideal.

Def: If I is an ideal in B , $I^* := \{-u : u \in I\}$ is called the dual filter of I .

Def: If F is a filter in B , then $F^* := \{-u : u \in F\}$ is called the dual ideal of F .

Def: The dual filter $I_0^* := \{x \subseteq \omega : \omega \setminus x \text{ is finite}\}$ of the Fréchet ideal I_0 on $\mathcal{P}(\omega)$ is called the Fréchet filter.

Definitions: Let I be an ideal in B and F a filter in B .

I is called:

- trivial if $I = \{0\}$
- principal if $\exists u \in B : I = \{v : v \leq u\}$
- prime if $\forall u \in B$, either $u \in I$ or $-u \in I$

F is called:

- trivial if $F = \{1\}$
- principal if $\exists u \in B : F = \{v : v \geq u\}$
- an ultrafilter if $\forall u \in B$, either $u \in F$ or $-u \in F$

Ideals and filters over ω :

We consider the Boolean algebra $(\mathcal{P}(\omega), \cup, \cap, \setminus, \emptyset, \omega)$. A set $I \subseteq \mathcal{P}(\omega)$ is an ideal over ω if $\forall x, y \subseteq \omega$:

- $\emptyset \in I$ and $\omega \notin I$
- $(x \in I \wedge y \in I) \Rightarrow x \cup y \in I$
- $(x \in I \vee y \in I) \Rightarrow x \cap y \in I$

A set F is a filter over ω if $\forall x, y \subseteq \omega$:

- $\emptyset \notin F$ and $\omega \in F$
- $(x \in F \wedge y \in F) \Rightarrow x \cap y \in F$
- $(x \in F \vee y \in F) \Rightarrow x \cup y \in F$

Remarks: The trivial ideal over ω is $\{\emptyset\}$ and the trivial filter over ω is $\{\omega\}$

- For any non-empty subset $x \subseteq \omega$, $F_x := \{y \in \mathcal{P}(\omega) : y \supseteq x\}$ is a principal filter and the dual ideal $I_{w_x} := (F_x)^* = \{z \in \mathcal{P}(\omega) : \omega \setminus z \in F_x\} = \{z \in \mathcal{P}(\omega) : z \cap x = \emptyset\}$ is a principal ideal. In particular, if $x = \{a\}$ for some $a \in \omega$, then F_x is a principal ultrafilter and I_{w_x} is a principal prime ideal.

Claim: Every principal ultrafilter over ω is of the form

$$F_{\{a\}} = \{ \gamma \in \mathcal{P}(\omega) : \{a\} \in \gamma \} \text{ for some } a \in \omega$$

Proof: Let F be a principal ultrafilter ^{over ω} . Because it is principal $\exists u \in \omega$ such that $F = \{ \gamma : u \in \gamma \}$. Take $\gamma \in F$. Because F is an ultrafilter either $\gamma \in F$ or $\mathcal{P}(\omega) \setminus \gamma \in F$. The first case is not possible, thus the second case must hold and then $\gamma \in F$ and consequently $u \in \gamma$. Thus, $\gamma \in F$ and consequently u consists of a singleton. $\square$

Remark: We denote that the Fréchet filter is neither a principal filter nor a ultrafilter.

Similarly, the Fréchet ideal is neither prime nor principal.

We consider now ultrafilters over arbitrary non-empty sets S .

Fact: Let $\mathcal{U}$ be an ultrafilter over S .

- a) If $\{x_0, \dots, x_{n-1}\} \subseteq \mathcal{P}(S)$ for some $n \in \omega$ such that $x_0 \cup \dots \cup x_{n-1} \in \mathcal{U}$ and for any distinct $i, j \in n$, we have $x_i \cap x_j \notin \mathcal{U}$, then there is a unique $k \in n$ such that $x_k \in \mathcal{U}$.
- b) If $x \in \mathcal{U}$ and $|x| \geq 2$, then there is a proper subset $y \subsetneq x$ such that $y \in \mathcal{U}$.
- c) If $\mathcal{U}$ contains a finite set, then $\mathcal{U}$ is principal.

Proof: a) Existence: Assume $x_k \notin \mathcal{U} \forall k \in n$. Then, $-x_k \in \mathcal{U} \forall k \in n$, which implies that

$$-x_0 \cap \dots \cap -x_{n-1} = -(x_0 \cup \dots \cup x_{n-1}) \in \mathcal{U} \Rightarrow x_0 \cup \dots \cup x_{n-1} \notin \mathcal{U} \text{ } \square$$

Uniqueness: Consider $i, j \in n$ distinct. Assume $x_i \in \mathcal{U}$ and $x_j \in \mathcal{U}$. Then $x_i \cap x_j \in \mathcal{U} \text{ } \square$

- b) If for any proper subset of x , $y \subsetneq x$, we have $y \notin \mathcal{U}$. Then $-y \in \mathcal{U} \Rightarrow \bigcap_{y \subsetneq x} -y = -\bigcup_{y \subsetneq x} y$
 $= -x \in \mathcal{U} \Rightarrow x \notin \mathcal{U} \text{ } \square$ Since $|x| \geq 2$, the existence of proper subsets in $\mathcal{U}$ is ensured.

- c) Let $A = \{a_1, \dots, a_n\}$ be a finite set in $\mathcal{U}$. By fact 1, we have that $\exists! k$ such that $\{a_k\} \in \mathcal{U}$. $F_{\{a_k\}} \subseteq \mathcal{U}$: Take $\gamma \in F_{\{a_k\}}$. Then $\gamma \cup \{a_k\} = \gamma \in \mathcal{U}$ by the third property of a filter. Since $\mathcal{U}$ is a ultrafilter, each $x \in -F_{\{a_k\}}$ does not belong to $\mathcal{U}$. Thus, $\mathcal{U} = F_{\{a_k\}}$ and $\mathcal{U}$ is principal.

Theorem: (Prime ideal theorem)

If I is an ideal in a Boolean algebra, then I can be extended to a prime ideal

Recall of A.S (26.03):

Definition: A family $\mathcal{F}$ of sets is said to have finite character if for each set x ,

$x \in \mathcal{F} \Leftrightarrow \text{fin}(x) \subseteq \mathcal{F}$, i.e. every finite subset of x belongs to $\mathcal{F}$

Techmüller's principle:

Let $\mathcal{F}$ be a non-empty family of sets.

If $\mathcal{F}$ has finite character, then $\mathcal{F}$ has a maximal element (with respect to inclusion " $\subseteq$ ")

Theorem 6.1.

Axiom of Choice $\Leftrightarrow$ Techmüller's principle

Proposition 6.6.

Axiom of Choice $\Rightarrow$ Principle ideal theorem

Proof:

- By Theorem 6.1, it suffices to prove:
Techmüller's principle $\Rightarrow$ Principle ideal theorem
- Let $(B, +, \cdot, -, 0, 1)$ be a Boolean algebra, let $I_0 \subsetneq B$ be an ideal
- Let $\mathcal{F}$ be the family of all sets $X \subseteq B \setminus I_0$ s.t. for every finite subset $\{v_0, \dots, v_n\} \subseteq X \cup I_0$, we have
$$v_0 + \dots + v_n \neq 1$$
- Take $X \in \mathcal{F}$, and $S \in \text{fin}(X)$, we know that
$$\forall A \in \text{fin}(X \cup I_0): v_0 + \dots + v_n \neq 1$$
$$\{v_0, \dots, v_n\}$$
Since $\text{fin}(S \cup I_0) \subseteq \text{fin}(X \cup I_0)$, it's clear that:

$$\forall A \in \text{fin}(S \cup I_0) : v_0 + \dots + v_n \neq 1$$

$$\Rightarrow S \in \mathcal{F}$$

Now, let $X \subseteq B$, assume that

$$\forall S \in \text{fin}(X) : S \in \mathcal{F}$$

Suppose now that $X \notin \mathcal{F}$, i.e.

$$1) X \not\subseteq B \setminus I_0 \quad \text{or} \quad 2) \exists \{v_0, \dots, v_n\} \in X \cup I_0 \text{ s.t. } v_0 + \dots + v_n \neq 1$$

1) cannot hold because otherwise,

$$\exists v \in X \text{ s.t. } v \notin X \Rightarrow \{v\} \notin \mathcal{F}, \text{ but } \{v\} \in \text{fin}(X) \nsubseteq \mathcal{F}$$

$$2) \text{ Take } U := \bigcup_{\substack{v_i \notin I_0 \\ 0 \leq i \leq n}} \{v_i\}, \text{ then } U \notin \mathcal{F} \text{ but } U \in \text{fin}(X) \nsubseteq \mathcal{F}$$

So, 2) cannot hold and we have, $X \in \mathcal{F} \Leftrightarrow \text{fin}(X) \in \mathcal{F}$

$\Rightarrow \mathcal{F}$ has finite character

^{Tell}
 $\Rightarrow \mathcal{F}$ has a maximal element I_1

^{Principle}
In particular, $I_1 \cap I_0 = \emptyset$

• Let $I := I_0 \cup I_1$, then for finitely many elements $\{v_0, \dots, v_n\} \in I$, we have that:

$$v_0 + \dots + v_n \neq 1 \quad (*)$$

• Claim 1, I is an ideal

$$\triangleright \overset{\text{Ideal}}{0} \in I_0 \subseteq I$$

$1 \notin I_0$ and $1 \notin I_1$, because otherwise, we would

obtain a contradiction to $(*) \Rightarrow 1 \notin I$

$\triangleright$ Let $u, v \in I$,

$\rightarrow$ If $u, v \in I_0$, then $u+v \in I_0 \subseteq I$

$\rightarrow$ If $u \in I_0, v \in I_1$,

take $X = \{u+v\}$, then $X \in \mathcal{F}$ because $X \not\subseteq \mathcal{F}$ would contradict that $v \in I_1$

So, $X \subseteq I_1$ by maximality of I_1

$$\Rightarrow u+v \in I_1 \subseteq I$$

→ If $u, v \in I_1$, again by maximality of I_1 , we have that $u+v \in I_1 \subseteq I$

▶ Let $w \in B$, $u \in I$

→ If $u \in I_0$, then $w \cdot u \in I_0 \subseteq I$

→ If $u \in I_1$, by absorption 4)

$$u + (u \cdot w) = u \quad (**)$$

Suppose $u+w \notin I_1$,

$$\stackrel{\text{max}}{\Rightarrow} \exists \{u_0 + \dots + u_n\} \in I_0 \cup I_1 : (u+w) + u_0 + \dots + u_n = 1$$

By iv), this implies.

$$u + ((u+w) + u_0 + \dots + u_n) = 1$$

$$\stackrel{\text{dist}}{\Rightarrow} u + u_0 + \dots + u_n = 1 \Rightarrow u \notin I_1 \quad \square$$

Hence, I is an ideal

• Claim 2. I is prime

Proof.

By maximality of I_1 ,

$$\forall v \in B \setminus I \exists u \in I : u+v=1$$

By uniqueness of the complement,

$$\forall v \in B \setminus I : -v \in I$$

So, for all u in B , either $u \in I$ or $-u \in I$

$\Rightarrow I$ is prime

• Hence, I is a prime ideal in B which extends I_0 $\square$

Ultrafilter theorem: If F is a filter over a set S , then F can be extended to an ultrafilter

We introduce the following definition:

Def: Let S be a set and $\mathcal{B}$ a set of binary functions (i.e. with values 0 or 1) defined on finite subsets of S . We say that $\mathcal{B}$ is a binary mess on S if $\mathcal{B}$ satisfies:

a) For each finite $P \subseteq S$, there is a function $g \in \mathcal{B}$ such that $\text{dom}(g) = P$, i.e. g is defined on P .

b) For each $g \in \mathcal{B}$ and each finite set $P \subseteq S$, $g|_P$ belongs to $\mathcal{B}$.

Def: Let f be a binary function on S and let $\mathcal{B}$ be a binary mess on S . We say that f is consistent with $\mathcal{B}$ if for every finite set $P \subseteq S$, we have $f|_P \in \mathcal{B}$.

We can now state the following:

Consistency Principle: For every binary mess $\mathcal{B}$ on a set S , there exists a binary function f on S which is consistent with $\mathcal{B}$.

We introduce now some terminology from Propositional Logic.

The alphabet of propositional logic consists of an arbitrarily large but fixed set

$\mathcal{P} := \{p_\lambda : \lambda \in \Lambda\}$ of so-called propositional variables, as well as of the logic operators " $\neg$ ", " $\wedge$ " and " $\vee$ ".

Definition: The formulae of propositional logic are defined recursively as follows:

- A single propositional variable $p \in \mathcal{P}$ by itself is a formula.

- If φ and ψ are formulae, then so are $\neg\varphi$, $\varphi \wedge \psi$ and $\varphi \vee \psi$.

Definition: A realisation of propositional logic is a map of $\mathcal{P}$, the set of propositional variables to the two element Boolean algebra $(\{0, 1\}, +, \cdot, -, 0, 1)$

Definition: Let φ be any formula. If a realisation f maps φ to 1, then we say that f satisfies φ . A set Σ of formulae of propositional logic is satisfiable if there is a realisation which simultaneously satisfies all the formulae in Σ .

Compactness theorem for propositional logic: Let Σ be a set of formulae of propositional logic. If every finite subset of Σ is satisfiable, then Σ is also satisfiable.

We can now state the following theorem:

Theorem: The following statements are equivalent:

- a) Prime ideal theorem
- b) Ultrafilter theorem
- c) Consistency principle
- d) Compactness theorem for propositional logic
- e) Every Boolean algebra has a prime ideal

Proof: ("a) $\Rightarrow$ b") The ultrafilter theorem is an immediate consequence of the dual form of the prime ideal theorem

"b) $\Rightarrow$ c)" Let $\mathcal{B}$ be a binary mess on a non-empty set S . We have to show that there is a binary function f on S which is consistent with $\mathcal{B}$, assuming that the ultrafilter theorem holds.

Let $\text{fin}(S)$ be the set of all finite subsets of S . For each $P \in \text{fin}(S)$, define

$\mathcal{X}_P := \{g \in {}^S 2 : g|_P \in \mathcal{B}\}$. Since $\mathcal{B}$ is a binary mess, we have that for each $g \in \mathcal{B}$ and
 $\uparrow$ set of binary functions over S

for each finite set P $g|_P \in \mathcal{B}$. Thus, the intersection of finitely many $\mathcal{X}_P$ is non-empty.

We can notice that the family $\mathcal{F}$ of all supersets of intersections of finitely many sets $\mathcal{X}_P$ is a filter over ${}^S 2$. Indeed,

- $\emptyset \notin \mathcal{F}$ and ${}^S 2 \in \mathcal{F}$,
- Take two supersets of intersections of finitely many $\mathcal{X}_P$. Their intersection is also a superset of intersections of finitely many $\mathcal{X}_P$.
- Take $A \in {}^S 2$ and $B \in \mathcal{F}$. Then, $A \cup B \in \mathcal{F}$ because it is a superset of B , which is itself a superset of intersections of finitely many $\mathcal{X}_P$.

Since $\mathcal{F}$ is a filter, we can extend it by the ultrafilter theorem to an ultrafilter $\mathcal{U} \in \mathcal{P}(\mathcal{S}_2)$. Since $\mathcal{U}$ is an ultrafilter, for each $x \in S$, either $\{g \in \mathcal{S}_2 : g(x) = 0\}$ or $\{g \in \mathcal{S}_2 : g(x) = 1\}$ belongs to $\mathcal{U}$. Define $f \in \mathcal{S}_2$ by stipulating that $\chi_{\{x\}} := \{g \in \mathcal{S}_2 : g(x) = f(x)\}$ belongs to $\mathcal{U}$. For any finite set $P = \{x_0, \dots, x_n\} \subseteq S$, we have that $\bigcap_{i \in n} \chi_{\{x_i\}} \in \mathcal{U}$, since $\mathcal{U}$ is a filter and thus $f|_P \in \mathcal{B}$, i.e. f is consistent with $\mathcal{B}$.

"c) $\Rightarrow$ d)" Let Σ be a set of formulae of propositional logic and let $S \subseteq \mathcal{P}$ be the set of propositional variables which appear in formulae of Σ . Assume that every finite subset is satisfiable, i.e. for every finite subset $\Sigma_0 \subseteq \Sigma$ there is a realisation $g_{\Sigma_0} : S_{\Sigma_0} \rightarrow \{0, 1\}$ which satisfies Σ_0 , where S_{Σ_0} denotes the set of propositional variables which appear in formulae of Σ_0 .

Assuming the consistency principle, we have to show that Σ is also satisfiable.

Define $\mathcal{B}_\Sigma := \{g_{\Sigma_0}|_P : \Sigma_0 \in \text{fin}(\Sigma) \wedge P \subseteq S_{\Sigma_0}\}$

We denote that: • For each $P \subseteq S$ finite, we have by assumption a g_{Σ_0} in $\mathcal{B}_\Sigma$ such that g_{Σ_0} is defined on P .

• For each $g \in \mathcal{B}_\Sigma$ and each finite $P \subseteq S$, $g|_P \in \mathcal{B}_\Sigma$.

Thus $\mathcal{B}_\Sigma$ is a binary nest. By the consistency principle, there exists a binary function f on S which is consistent with $\mathcal{B}_\Sigma$. By the definition of $\mathcal{B}_\Sigma$, it means that f is a realisation which satisfies Σ .

d) $\Rightarrow$ e).

- Let $(B, +, \cdot, -, 0, 1)$ be a Boolean algebra. Let $P_i = \{p_i \mid i \in B\}$, a set of propositional variables. Let Σ_B be the following set of formulae of Propositional Logic.

I). $p_0, \neg p_1$

II). $p_0 \vee \neg p_0$ (for each $i \in B$)

III). $\neg (p_{i_1} \wedge \dots \wedge p_{i_n}) \vee p_{i_1 + \dots + i_n}$ (for each finite set $\{i_1, \dots, i_n\} \in B$)

IV). $\neg (p_{i_1} \vee \dots \vee p_{i_n}) \vee p_{i_1 - i_n}$ (for each finite set $\{i_1, \dots, i_n\} \in B$)

- Every finite subset of B generates a finite subalgebra of B and every finite Boolean algebra has a prime ideal.
- Every finite prime ideal in a finite subalgebra of B corresponds to a finite subset of Σ_B , and vice versa $\stackrel{d)}{\Rightarrow} \Sigma_B$ is satisfiable.
- Let f be a realisation of Σ_B , define $I = \{i \in B \mid f(p_i) = 1\} \subseteq B$.
- Claim 1: I is an ideal.

Proof:

$$\Rightarrow f(p_0) \stackrel{\text{I)}}{=} 1 \text{ and } -f(p_1) = f(\neg p_1) \stackrel{\text{I)}}{=} 1 \stackrel{\text{vi)}}{\Rightarrow} f(p_1) = 0 \\ \Rightarrow p_0 \in I, p_1 \notin I$$

$$\Rightarrow \text{If } f(p_{i_1}) = f(p_{i_2}) = 1, \text{ then}$$

$$\begin{aligned} 1 &\stackrel{\text{III)}}{=} f(\neg (p_{i_1} \wedge p_{i_2}) \vee p_{i_1 + i_2}) \\ &= f(\neg (p_{i_1} \wedge p_{i_2})) + f(p_{i_1 + i_2}) \\ &= -f(p_{i_1} \wedge p_{i_2}) + f(p_{i_1 + i_2}) \\ &= -(f(p_{i_1}) \cdot f(p_{i_2})) + f(p_{i_1 + i_2}) \\ &\stackrel{\text{De Morgan}}{=} -f(p_{i_1}) + -f(p_{i_2}) + f(p_{i_1 + i_2}) \\ &\stackrel{\text{vii)}}{=} 0 + 0 + 1 = 1 \end{aligned}$$

$$= f(p_{v_1+v_2})$$

$$\Rightarrow \forall v_1, v_2 \in I, v_1+v_2 \in I$$

→ Analogously, let $v_1 \in I, v_2 \in B$

$$\text{IV)} \Rightarrow f(p_{v_2 \cdot v_1}) = 1 \Rightarrow v_2 \cdot v_1 \in I$$

• Claim 2, I is prime

Proof, Suppose $v \notin I$

By II,

$$f(p_v) = 0 \Rightarrow p_v \notin \sum_B \stackrel{\text{II}}{\Rightarrow} f(\neg p_v) = 1$$

If $v \in I$,

$$f(\neg p_v) = -f(p_v) = -1 = 0$$

□

e) $\Rightarrow$ a).

• Let $(B, +, \cdot, -, 0, 1)$ be a Boolean algebra and $I \subseteq B$ be an ideal

Goal, Show that I is contained in some prime ideal in B

• Define an equivalence relation on B .

$$u \sim v \Leftrightarrow (u \cdot (-v)) + (v \cdot (-u)) \in I$$

Notice that this means:

$$u \sim v \Leftrightarrow u \oplus v \in I$$

(Quick check that this is indeed an equivalence relation.)

$$\rightarrow u \sim u \text{ because } u \cdot (-u) = 0 \in I$$

$$\rightarrow u \sim v \Leftrightarrow v \sim u, \text{ clear by def.}$$

$$\rightarrow u \sim v, v \sim z, \text{ then}$$

$$u \oplus z = u \oplus \underbrace{v \oplus v}_{=0} \oplus z$$

$$= \underbrace{(u \oplus v)}_{\in I} \cdot (- (v \oplus z)) + (- (u \oplus v) \cdot \underbrace{(v \oplus z)}_{\in I}) \in I$$

• Let $C := B/\sim$, with $\oplus$ on C .

$$[u] + [v] = [u+v]$$

$$[u] \cdot [v] = [u \cdot v]$$

$$-[u] = [-u]$$

Then, $(C, +, \cdot, -, [0], [1])$ is a Boolean algebra, called the quotient of B modulo I

• e) $\Rightarrow C$ has a prime ideal $\mathcal{J} \subseteq C$

• Claim: $P := \{u \in B \mid [u] \in \mathcal{J}\}$ is a prime ideal in B which extends I

Part 1. $I \subseteq P$

Let $u \in I$, then $u \sim 0$ because

$$\begin{aligned} (u \cdot (-u)) + (u \cdot (-u)) &\stackrel{\text{vi) + iii)}}{=} u \cdot 1 + 0 \\ &\stackrel{\text{v)}}{=} u + 0 \\ &\stackrel{\text{ii)}}{=} u \in I \end{aligned}$$

$$[0] \in \mathcal{J}, [0] = [u] \Rightarrow [u] \in \mathcal{J} \Rightarrow u \in P$$

Part 2. P is a prime ideal

$\Rightarrow 0 \in P$ because $[0] \in \mathcal{J}$

$1 \notin P$ because $[1] \notin \mathcal{J}$

$\Rightarrow$ Let $u, v \in P$, then

$$[u+v] = [u] + [v] \in \mathcal{J} \Rightarrow u+v \in P$$

$\Rightarrow$ Let $w \in B, u \in P$, then

$$[w \cdot u] = [w] \cdot [u] \in \mathcal{J} \Rightarrow w \cdot u \in P$$

So actually,

$\mathcal{J}$ is an ideal $\Rightarrow P$ is an ideal

$\Rightarrow$ Let $u \in B$,

if $u \notin P$, then $[u] \notin \mathcal{J}$

$\mathcal{J}$ prime $\Rightarrow -[u] = [-u] \in \mathcal{J} \Rightarrow -u \in P$

Analogously, if $u \in P$, then $-u \notin P$

So we have.

$\mathcal{J}$ prime $\Rightarrow P$ prime

