

On the Largest Eigenvalue of a Random Subgraph of the Hypercube

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Abstract: Let G be a random subgraph of the n -cube where each edge appears randomly and independently with probability p . We prove that the largest eigenvalue of the adjacency matrix of G is almost surely

$$\lambda_1(G) = (1 + o(1)) \max(\Delta^{1/2}(G), np),$$

where $\Delta(G)$ is the maximum degree of G and the $o(1)$ term tends to zero as $\max(\Delta^{1/2}(G), np)$ tends to infinity.

1. Introduction and Formulation of Results

Let Q^n be a graph whose vertices are all the vectors $\{x = (x_1, \dots, x_n) \mid x_i \in \{0, 1\}\}$ and two vectors x and y are adjacent if they differ in exactly one coordinate, i.e., $\sum_i |x_i - y_i| = 1$. We call Q^n the n -dimensional cube or simply the n -cube. Clearly Q^n is an n -regular, bipartite graph of order 2^n . In this paper we study random subgraphs of the n -cube. A random subgraph $G(Q^n, p)$ is a discrete probability space composed of all subgraphs of the n -cube, where each edge of Q^n appears randomly and independently with probability $p = p(n)$. Sometimes with some abuse of notation we will refer to a random subgraph $G(Q^n, p)$ as a graph on 2^n vertices generated according to the distribution described above. Usually asymptotic properties of random graphs are of interest. We say that a graph property \mathcal{A} holds *almost surely*, or a.s. for brevity, in $G(Q^n, p)$ if the probability that $G(Q^n, p)$ has property \mathcal{A} tends to one as n tends to infinity. Necessary background information on random graphs in general and random subgraphs of n -cube can be found in [4].

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Random subgraphs of the hypercube were studied by Burtin [5], Erdős and Spencer [8], Ajtai, Komlós and Szemerédi [2] and Bollobás [4], among others. In particular it was shown that a giant component emerges shortly after $p = 1/n$ ([2]) and the graph becomes connected shortly after $p = 1/2$ ([5, 8, 4]). Recently the model has become of interest in mathematical biology ([7, 13, 14]). In particular it appears (see [13, 14]) that random graphs play an important role in a general model of a population evolving over a network of selectively neutral genotypes. It has been shown that the population's limit distribution on the neutral network is solely determined by the network topology and given by the principal eigenvector of the network's adjacency matrix. Moreover, the average number of neutral mutant neighbors per individual is given by the spectral radius.

The subject of this paper is the asymptotic behavior of the largest eigenvalue of the random graph $G(Q^n, p)$. The adjacency matrix of G is an $2^n \times 2^n$ matrix A whose entries are either one or zero depending on whether the edge (x, y) is present in G or not. A is a random real symmetric matrix with the eigenvalues denoted by $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_{2^n}$. It follows from the Perron-Frobenius theorem that the largest eigenvalue is equal to the spectral norm of A , i.e. $\lambda_{\max}(G) = \lambda_1(G) = \|A\| = \max_j |\lambda_j|$. Note also that for a subgraph of the n -cube, or in general, for any bipartite graph, $\lambda_k(G) = -\lambda_{|V|-k}(G)$, $k = 1, 2, \dots$ and in particular $|\lambda_{\min}(G)| = \lambda_{\max}(G)$. It is easy to observe that for every graph $G = (V, E)$ its largest eigenvalue $\lambda_1(G)$ is always squeezed between the average degree of G , $\bar{d} = \sum_{v \in V} d_G(v)/|V|$ and its maximal degree $\Delta(G) = \max_{v \in V} d_G(v)$. In some situations these two bounds have the same asymptotic value which determines the behavior of the largest eigenvalue. For example, this easily gives the asymptotics of the largest eigenvalue of a random subgraph $G(n, p)$ of a complete graph of order n for $p \gg \log n/n$. On the other hand, in our case there is a gap between average and maximal degree of random graph $G(Q^n, p)$ for all values of $p < 1$ and therefore it is not immediately clear how to estimate its largest eigenvalue.

Here we determine the asymptotic value of the largest eigenvalue of sparse random subgraphs of n -cube. To understand better the result, observe that if Δ denotes a maximal degree of a graph G , then G contains a star S_Δ and therefore $\lambda_1(G) \geq \lambda_1(S_\Delta) = \sqrt{\Delta}$. Also, as mentioned above $\lambda_1(G)$ is at least as large as the average degree of G . As for all values of $p(n) \gg n^{-1}2^{-n}$, a.s. $|E(G(Q^n, p))| = (1 + o(1))pn2^n$, we get that a.s. $\lambda_1(G(Q^n, p)) \geq (1 + o(1))np$. Combining the above lower bounds, we get that a.s. $\lambda_1(G(Q^n, p)) \geq (1 + o(1)) \max(\sqrt{\Delta}, np)$. It turns out this lower bound can be matched by an upper bound of the same asymptotic value, as given by the following theorem:

Theorem 1.1. *Let $G(Q^n, p)$ be a random subgraph of the n -cube and let Δ be the maximum degree of G . Then almost surely the largest eigenvalue of the adjacency matrix of G satisfies*

$$\lambda_1(G) = (1 + o(1)) \max(\sqrt{\Delta}, np),$$

where the $o(1)$ term tends to zero as $\max(\Delta^{1/2}(G), np)$ tends to infinity.

As the asymptotic value of the maximal degree of $G(Q^n, p)$ can be easily determined for all values of $p(n)$ (see Lemma 2.1), the above theorem enables us to estimate the asymptotic value of $\lambda_1(G(Q^n, p))$ for all relevant values of p . This theorem is similar to the recent result of Krivelevich and Sudakov [11] on the largest eigenvalue of a random subgraph $G(n, p)$ of a complete graph of order n .

The rest of this paper is organized as follows. In the next section we gather some necessary technical lemmas about random subgraphs of the n -cube. Section 3 is devoted

to the proof of the main theorem. Section 4, the last section of the paper, contains some concluding remarks.

Throughout the paper we will systematically omit floor and ceiling signs for the sake of clarity of presentation. All logarithms are natural. We will frequently use the inequality $\binom{a}{b} \leq (ea/b)^b$. Also we use the following standard notations: $a_n = \Theta(b_n)$, $a_n = O(b_n)$ and $a_n = \Omega(b_n)$ for $a_n > 0$, $b_n > 0$ as $n \rightarrow \infty$ if there exist constants C_1 and C_2 such that $C_1 b_n < a_n < C_2 b_n$, $a_n < C_2 b_n$ or $a_n > C_1 b_n$ respectively. The equivalent notations $a_n = o(b_n)$ and $a_n \ll b_n$ mean that $a_n/b_n \rightarrow 0$ as $n \rightarrow \infty$. We will say that an event Υ_n , depending on a parameter n , holds almost surely if $\Pr(\Upsilon_n) \rightarrow 1$ as $n \rightarrow \infty$ (please note that this notion, very common in the literature on random structures, has a different meaning in probability theory).

2. Few Technical Lemmas

In this section we establish some properties of random subgraphs of the n -cube which we will use later in the proof of our main theorem. First we consider the asymptotic behavior of the maximal degree of $G(Q^n, p)$. It is not difficult to show that if p is a constant less than $1/2$ then a.s. $\Delta(G) = (1 + o(1))cn$, where c satisfies the equation $\log 2 + c \log p + (1-c) \log(1-p) = c \log c + (1-c) \log(1-c)$ and $\Delta(G) = (1+o(1))n$ for $p \geq 1/2$. We omit the proof of this statement since for our purposes it is enough to have $\sqrt{\Delta(G)} = O(np)$ which follows immediately from the fact that $\Delta(G) \leq n$. The case when $p = o(1)$ is studied in more detail in the following lemma.

Lemma 2.1. *Let $G = G(Q^n, p)$ be a random subgraph of the n -cube. Denote by*

$$\kappa(n) = \max \left\{ k : 2^n \binom{n}{k} p^k (1-p)^{n-k} \geq 1 \right\}.$$

Then the following statements hold.

- (i) *If $p = o(1)$ and p is not at least exponentially small in n then almost surely $\kappa(n) - 1 \leq \Delta(G) \leq \kappa(n) + 1$.*
- (ii) *If $p = \Theta(2^{-n/k} n^{-1})$, then $2^n \binom{n}{k} p^k (1-p)^{n-k} = \Theta(1)$ and $\kappa(n) = k - 1$ or k . Also almost surely $\Delta(G)$ is either $k - 1$ or k .*
- (iii) *If p is exponentially small, but not proportional to $2^{-n/k} n^{-1}$, then $\kappa(n) = \left\lceil \frac{n \log 2}{\log(p^{-1}) - \log n} \right\rceil$ and almost surely $\Delta(G) = \kappa(n)$.*

Proof. Let X_k be the number of vertices of $G(Q^n, p)$ with degree larger than $k - 1$. Then $X_k = \sum_{v \in Q^n} I_v$, where I_v is the indicator random variable of the event that $\deg(v) \geq k$.

One can easily calculate the expectation $\mathbf{E}X_k = 2^n \sum_{l \geq k} \binom{n}{l} p^l (1-p)^{n-l}$. Also note that Q^n is bipartite and therefore has an independent set X of size 2^{n-1} . By definition, the events that $d(v) < k$ are mutually independent for all $v \in X$. Therefore we obtain that

$$\begin{aligned} \Pr(X_k = 0) &\leq \prod_{v \in X} \Pr(d(v) < k) = \prod_{v \in X} (1 - \mathbf{E}I_v) \\ &\leq \exp\left(-\sum_{v \in X} \mathbf{E}I_v\right) = \exp(-\mathbf{E}X_k/2). \end{aligned} \tag{1}$$

Let us now consider the case (i) in more detail. We have that

$$2^{n-O(np)} (n/k)^k p^k \leq 2^n \binom{n}{k} p^k (1-p)^{n-k} \leq 2^n (en/k)^k p^k.$$

Therefore it is easy to check that $\kappa(n)$ must satisfy the inequalities

$$\frac{n \log 2(1 - 1/\log \log(p^{-1}))}{\log(p^{-1})} \leq \kappa(n) \leq \frac{n \log 2(1 + 1/\log \log(p^{-1}))}{\log(p^{-1})}.$$

By definition, $\mathbf{E}X_{\kappa(n)} \geq 1$. Elementary computations show that $\mathbf{E}X_{k+1} = (1 + o(1)) \frac{p \log(p^{-1})}{\log 2} \mathbf{E}X_k$ for $k = (1 + o(1))\kappa(n)$ which imply $\mathbf{E}X_{\kappa(n)-1} \geq (1 + o(1)) \frac{\log 2}{p \log(p^{-1})}$. Therefore by (1), we have that

$$\begin{aligned} \Pr(\Delta(G) < \kappa(n) - 1) &= \Pr(X_{\kappa(n)-1} = 0) \leq \exp\left(-\frac{\mathbf{E}X_{\kappa(n)-1}}{2}\right) \\ &\leq \exp\left(-(1 + o(1)) \frac{\log 2}{p \log(p^{-1})}\right) = o(1). \end{aligned}$$

On the other hand, since $\mathbf{E}X_{\kappa(n)+1} \leq 1 + o(1)$ we have that $\mathbf{E}X_{\kappa(n)+2} \leq (1 + o(1)) \frac{p \log(p^{-1})}{\log 2} = o(1)$. Thus by Markov's inequality we conclude that a.s. $X_{\kappa(n)+2} = 0$ and therefore almost surely $\kappa(n) - 1 \leq \Delta(G) \leq \kappa(n) + 1$.

Now consider the case (ii). Since $p(n) = \Theta(2^{-n/k} n^{-1})$ we have that $2^n \binom{n}{k} p^k (1-p)^{n-k} = \Theta(1)$ and $\kappa(n) = k-1$ or k . Also it is easy to check that $\mathbf{E}X_{k+1} = \Theta(2^{-n/k}) = o(1)$ and $\mathbf{E}X_{k-1} = \Theta(2^{n/k})$. Therefore by (1) we have that $\Pr(\Delta(G) < k-1) \leq \exp(-\mathbf{E}X_{k-1}/2) = o(1)$ and by Markov's inequality $\Pr(\Delta(G) \geq k+1) = o(1)$.

Finally suppose that $p(n)$ is exponentially small but not proportional to $2^{-n/k} n^{-1}$ for any $k \geq 1$. Then it is rather straightforward to check that $\kappa(n) = \left\lceil \frac{n \log 2}{\log(p^{-1}) - \log n} \right\rceil$ and $\mathbf{E}X_{\kappa(n)+1} \ll 1 \ll \mathbf{E}X_{\kappa(n)}$. Therefore, again using (1) and Markov's inequality, we conclude that $\Pr(\Delta(G) < \kappa(n)) = o(1)$ and $\Pr(\Delta(G) \geq \kappa(n) + 1) = o(1)$. Thus almost surely $\Delta(G) = \kappa(n)$. This completes the proof.

Next we need the following lemma, which shows that a.s. G cannot have a large number of vertices of high degree too close to one another. More precisely the following is true.

Lemma 2.2. *Let $G(Q^n, p)$ be a random subgraph of n -cube. Then almost surely*

- (i) *For every $0 < p \leq 1$ and for any two positive constants a and b such that $a + b > 1$ and $n^b \geq 6np$, G contains no vertex which has within distance one or two at least n^a vertices of G with degree $\geq n^b$.*
- (ii) *For $p \geq n^{-2/3}$ and any constant $a > 0$, G contains no vertex which has within distance one or two at least n^a/p vertices of G with degree $\geq np + np/\log n$.*

Proof. We prove the lemma for the case of vertices of distance two, the case of vertices of distance one can be treated similarly. Note that since the n -cube is a bipartite graph the vertices which are within the same distance from a given vertex in Q^n are not adjacent.

- (i) Let X be the number of vertices of G which violate condition (i). To prove the statement we estimate the expectation of X . Clearly we can choose a vertex v of the n -cube in 2^n ways. Since there are at most n^2 vertices in Q^n within distance two from v we have at most $\binom{n^2}{n^a}$ possibilities to choose a subset S of size n^a of vertices which will have degree at least n^b . The probability that the degree of some vertex in S is at least n^b is bounded by $\binom{n}{n^b} p^{n^b}$. Note that these events are mutually independent, since S contains no edges of the n -cube. Therefore, using that $a + b > 1$ and $b > 0$, we obtain that

$$\begin{aligned} \mathbf{E}(X) &\leq 2^n \binom{n^2}{n^a} \left(\binom{n}{n^b} p^{n^b} \right)^{n^a} \leq 2^n n^{2n^a} \left(\frac{enp}{n^b} \right)^{n^a} \\ &\leq 2^{n+2n^a \log_2 n} 2^{-n^{a+b}} = o(1). \end{aligned}$$

Thus by Markov's inequality we conclude that almost surely no vertex violates condition (i).

- (ii) Let again X be the number of vertices of G which violate condition (ii). Similarly as before we have 2^n choices for vertex v and at most $\binom{n^2}{n^a/p}$ choices for the set S of vertices within distance two from v which will have degree $\geq np + np/\log n$. Since for all vertices $v \in G$ the degree $d(v)$ is binomially distributed with parameters n and p , then by a standard large deviation inequality (cf., e.g., [1], Appendix A)

$$\begin{aligned} \Pr \left[d(v) \geq t = np + np/\log n \right] &\leq e^{-(t-np)^2/2np + (t-np)^3/2(np)^2} \\ &= e^{-(1+o(1))np/(2\log^2 n)}. \end{aligned}$$

As we already mentioned, the events that vertices in S have degree $\geq np + np/\log n$ are mutually independent. Therefore, using that $p \geq n^{-2/3}$ and $a > 0$, we conclude that

$$\begin{aligned} \mathbf{E}(X) &\leq 2^n \binom{n^2}{n^a/p} \left(e^{-(1+o(1))np/(2\log^2 n)} \right)^{n^a/p} \\ &\leq 2^n n^{2n^a/p} e^{-(1+o(1))n^{1+a}/(2\log^2 n)} \\ &\leq 2^n e^{2n^{a+2/3} \log n} e^{-(1+o(1))n^{1+a}/(2\log^2 n)} = o(1). \end{aligned}$$

Thus we can complete the proof of the lemma using again Markov's inequality.

3. Proof of the Theorem

In this section we present our main result. We start by listing some simple properties of the largest eigenvalue of a graph, that we will use later in the proof. Most of these easy statements can be found in Chapter 11 of the book of Lovász [12].

Proposition 3.1. *Let G be a graph on n vertices and m edges and with maximum degree Δ . Let $\lambda_1(G)$ be the largest eigenvalue of the adjacency matrix of G . Then it has the following properties:*

- (I) $\max(\sqrt{\Delta}, \frac{2m}{n}) \leq \lambda_1(G) \leq \Delta$.
- (II) If $E(G) = \cup_i E(G_i)$ then $\lambda_1(G) \leq \sum_i \lambda_1(G_i)$. If in addition graphs G_i are vertex disjoint, then $\lambda_1(G) = \max_i \lambda_1(G_i)$.
- (III) If G is a bipartite graph then $\lambda_1(G) \leq \sqrt{m}$. Moreover if it is a star of size Δ then $\lambda_1(G) = \sqrt{\Delta}$.
- (IV) If the degrees on both sides of bipartition are bounded by Δ_1 and Δ_2 respectively, then $\lambda_1(G) \leq \sqrt{\Delta_1 \Delta_2}$.
- (V) For every vertex v of G let $W_2(G, v)$ denote the number of walks of length two in G starting at v . Then $\lambda_1(G) \leq \sqrt{\max_v W_2(G, v)}$. \square

Proof of the Theorem 1.1. We already derived in the introduction the lower bound of this theorem so we will concentrate on proving an upper bound. We will frequently use the following simple fact that between any two distinct vertices of the n -cube there are at most two paths of length two. We divide the proof into a few cases with respect to the value of the edge probability p . In each case we partition G into smaller subgraphs, whose largest eigenvalue is easier to estimate. We start with a rather easy case when the random graph is relatively sparse.

Case 1. Let $e^{-\log^4 n} \leq p \leq n^{-2/3}$. For these values of p , by Lemma 2.1, we have that $\Delta(G) \geq \Omega(\frac{n}{\log^4 n})$. Partition the vertex set of G into three subsets as follows. Let V_1 be the set of vertices with degree at most $n^{2/5}$, let V_2 be the set of vertices with degree larger than $n^{2/5}$ but smaller than $n^{4/7}$ and let V_3 be the set of vertices with degree at least $n^{4/7}$. Also let G_1 be a subgraph of G induced by V_1 , let G_2 be a subgraph induced by $V_2 \cup V_3$, let G_3 be a bipartite graph containing edges of G connecting V_1 and V_2 and finally let G_4 be a bipartite graph containing edges connecting V_1 and V_3 . By definition $G = \cup_i G_i$ and thus by claim (II) of Proposition 3.1 we obtain that $\lambda_1(G) \leq \sum_{i=1}^4 \lambda_1(G_i)$.

Since the maximum degree of graph G_1 is at most $n^{2/5}$, then by (I) it follows that $\lambda_1(G_1) \leq n^{2/5}$. The degrees of vertices of bipartite graph G_3 are bounded on one side by $n^{2/5}$ and on another by $n^{4/7}$. Hence using (IV) we conclude that $\lambda_1(G_3) \leq \sqrt{n^{2/5} n^{4/7}} = n^{17/35}$. Let v be a vertex of G_2 and let $N_2(G_2, v)$ be the set of vertices of G_2 which are within distance exactly two from v . Since between any two distinct vertices of \mathcal{Q}^n there are at most two paths of length two it is easy to see that the number of walks of length two in G_2 starting at v is bounded by $d_{G_2}(v) + 2|N_2(G_2, v)|$. Since every vertex of $V_2 \cup V_3$ has degree in G at least $n^{2/5}$, using Lemma 2.2 (i) with $a = 4/5$ and $b = 2/5$ we obtain that almost surely both $d_{G_2}(v)$ and $|N_2(G_2, v)|$ are bounded by $n^{4/5}$. Therefore for every vertex v in G_2 we have $W_2(G_2, v) \leq d_{G_2}(v) + 2|N_2(G_2, v)| \leq 3n^{4/5}$ and hence by (V) $\lambda_1(G_2) \leq \sqrt{3n^{4/5}} = \sqrt{3}n^{2/5}$.

Finally we need to estimate $\lambda_1(G_4)$. To do so consider partition of V_1 into two parts. Let V'_1 be the set of vertices in V_1 with at least two neighbors in V_3 and let $V''_1 = V_1 - V'_1$. Let G'_4 and G''_4 be bipartite graphs with parts (V'_1, V_3) and (V''_1, V_3) respectively. By definition, $G_4 = G'_4 \cup G''_4$ and hence $\lambda_1(G_4) \leq \lambda_1(G'_4) + \lambda_1(G''_4)$. Since the vertices in V''_1 have at most one neighbor in V_3 and the graph G''_4 is bipartite it follows that G''_4 is the union of vertex disjoint stars of size at most $\Delta(G)$. So by (III) we get $\lambda_1(G''_4) \leq \sqrt{\Delta(G)}$. Now let u be the vertex of V_3 with at least $2n^{1/2}$ neighbors in V'_1 . By definition, every neighbor of u in V'_1 has also an additional neighbor in V_3 , which is distinct from u . Therefore we obtain that there are at least $2n^{1/2}$ simple paths of length two from u to the set V_3 . Since between any two distinct vertices of the n -cube there are at most two paths of length two we obtain that u has at least $n^{1/2}$ other vertices of V_3 within distance

two. Since the degree of all these vertices is at least $n^{4/7}$ it follows from Lemma 2.2 (i) with $a = 1/2$ and $b = 4/7$ that a.s. there is no vertex u with this property. Therefore the degree of every vertex from V_3 in bipartite graph G'_4 is bounded by $2n^{1/2}$ and we also have that the degree of every vertex from V'_1 is at most $n^{2/5}$. So using again **(IV)** we obtain $\lambda_1(G'_4) \leq \sqrt{2n^{1/2}n^{2/5}} = \sqrt{2}n^{9/20}$. This implies the desired bound on $\lambda_1(G)$, since

$$\begin{aligned} \lambda_1(G) &\leq \sum_i \lambda_1(G_i) \leq n^{2/5} + \sqrt{3}n^{2/5} + n^{17/35} + \left(\sqrt{2}n^{9/20} + \sqrt{\Delta(G)} \right) \\ &= (1 + o(1))\sqrt{\Delta(G)}. \end{aligned}$$

Case 2. Let $p \geq n^{-4/9}$. This case when the random graph is dense is also quite simple. Indeed, partition the vertices of G into two parts. Let V_1 be the set of vertices with degree larger than $np + np/\log n$ and let V_2 be the rest of the vertices. Clearly $G = \cup_i G_i$, where G_1 is a subgraph induced by V_1 , G_2 is a subgraph induced by V_2 , and G_3 is a bipartite subgraph with bipartition (V_1, V_2) . By definition, the maximum degree of G_2 is at most $np + np/\log n$, implying $\lambda_1(G_2) \leq np + np/\log n$. Since every vertex in V_1 has degree at least $np + np/\log n$, by Lemma 2.2 (ii) with $a = 1/18$ we obtain that almost surely no vertex in G can have more than $n^a/p \leq n^{1/2}$ vertices in V_1 within distance one or two. In particular, this implies that the maximum degree in G_1 is at most $n^{1/2}$ and so $\lambda_1(G_1) \leq n^{1/2}$.

Partition V_2 into two parts. Let V'_2 be the set of vertices in V_2 with at least two neighbors in V_1 and let $V''_2 = V_2 - V'_2$. Let G'_3 and G''_3 be bipartite graphs with parts (V_1, V'_2) and (V_1, V''_2) respectively. By definition, $G_3 = G'_3 \cup G''_3$ and thus $\lambda_1(G_3) \leq \lambda_1(G'_3) + \lambda_1(G''_3)$. Since the vertices in V''_2 have at most one neighbor in V_1 and the graph G''_3 is bipartite it follows that G''_3 is the union of vertex disjoint stars of size at most $\Delta(G)$. So by **(III)** we get $\lambda_1(G''_3) \leq \sqrt{\Delta(G)} \leq n^{1/2}$. Now let u be the vertex of V_1 with at least $2n^{1/2}$ neighbors in V'_2 . By definition, every neighbor of u in V'_2 has also an additional neighbor in V_1 , which is distinct from u . Therefore we obtain that there are at least $2n^{1/2}$ simple paths of length two from u to the set V_1 . Since between any two distinct vertices there are at most two paths of length two we obtain that u has at least $n^{1/2}$ other vertices of V_1 within distance two. As we already explain in the previous paragraph, this almost surely does not happen. Thus the degree of every vertex from V_1 in bipartite subgraph G'_3 is bounded by $2n^{1/2}$ and we also have that the degree of every vertex from V'_2 is at most $np + np/\log n$. So using **(IV)** we obtain $\lambda_1(G'_3) \leq \sqrt{2n^{1/2}(np + np/\log n)}$. Now since $np \geq n^{5/9}$, it follows that all $\lambda_1(G_1), \lambda_1(G'_3), \lambda_1(G''_3) = o(np)$ and therefore

$$\lambda_1(G) \leq \sum_i \lambda_1(G_i) \leq \lambda_1(G_2) + o(np) \leq np + np/\log n + o(np) = (1 + o(1))np.$$

Case 3. Let $n^{-2/3} \leq p \leq n^{-4/9}$. This part of the proof is slightly more involved than two previous ones since in particular it needs to deal with a delicate case when np and $\sqrt{\Delta(G)}$ are nearly equal.

Partition the vertex set of G into four parts. Let V_1 be the set of vertices with degree at least $n^{2/3}$ and let V_2 be the set of vertices with degrees larger than $np + np/\log n$ but less than $n^{2/3}$. Let V_4 contain all vertices G which have at least one neighbor in V_1 and degree at most $np + np/\log n$. Finally let V_3 be the set of remaining vertices of G . Note

that by definition there are no edges between V_1 and V_3 and every vertex from V_3 also has degree at most $np + np/\log n$ in G .

We consider the following subgraphs of G . Let G_1 be the bipartite subgraph containing all the edges between V_1 and V_4 . Partition V_4 into two parts. Let V'_4 be the set of vertices in V_4 with at least two neighbors in V_1 and let $V''_4 = V_4 - V'_4$. Let G'_1 and G''_1 be bipartite graphs with parts (V_1, V'_4) and (V_1, V''_4) respectively. By definition, $G_1 = G'_1 \cup G''_1$ and thus $\lambda_1(G_1) \leq \lambda_1(G'_1) + \lambda_1(G''_1)$. Since the vertices in V''_4 have at most one neighbor in V_1 and the graph G''_1 is bipartite it follows that G''_1 is the union of vertex disjoint stars of size at most $\Delta(G)$. So by **(III)** we get $\lambda_1(G''_1) \leq \sqrt{\Delta(G)}$. Now let u be the vertex of V_1 with at least $2n^{2/5}$ neighbors in V'_4 . By definition, every neighbor of u in V'_4 has also an additional neighbor in V_1 , which is distinct from u . Therefore we obtain that there are at least $2n^{2/5}$ simple paths of length two from u to the set V_1 . Similar as before, this implies that u has at least $n^{2/5}$ other vertices of V_1 within distance two. By Lemma 2.2 (i) with $a = 2/5$ and $b = 2/3$ this almost surely does not happen. Thus the degree of every vertex from V_1 in bipartite subgraph G'_1 is bounded by $2n^{2/5}$ and we also have that the degree of every vertex from V'_4 is at most $np + np/\log n \leq n^{5/9}$. So using **(IV)** we obtain $\lambda_1(G'_1) \leq \sqrt{2n^{2/5}n^{5/9}} = \sqrt{2}n^{43/90} = o(\sqrt{\Delta})$. Therefore $\lambda_1(G_1) \leq \lambda_1(G'_1) + \lambda_1(G''_1) \leq (1 + o(1))\sqrt{\Delta}$.

Our second subgraph G_2 is induced by the set V_3 . By definition, the maximum degree in it is at most $np + np/\log n$ and therefore $\lambda_1(G_2) \leq (1 + o(1))np$. Crucially this graph is a vertex disjoint form G_1 which implies by **(II)** that

$$\lambda_1(G_1 \cup G_2) = \max(\lambda_1(G_1), \lambda_1(G_2)) \leq (1 + o(1)) \max(\sqrt{\Delta}, np).$$

Next we define the remaining graphs whose union with G_1 and G_2 equals to G and show that their largest eigenvalues contribute only smaller order terms in the upper bound on $\lambda_1(G)$. Let G_3 be the subgraph of G induced by the set $V_1 \cup V_2$. By definition, every vertex in G_3 has at least $np + np/\log n$ neighbors in G . Therefore by Lemma 2.2 (ii) with $a = 1/12$ we obtain that for every $v \in G_3$ there at most $n^a/p \leq n^{3/4}$ other vertices of G_3 within distance one or two. This implies that $d_{G_3}(v)$ and $|N_2(G_3, v)|$ are both bounded by $n^{3/4}$. Then, as we already show in Case 1, the total number of walks of length two starting at v is bounded by $d_{G_3}(v) + 2|N_2(G_3, v)| \leq 3n^{3/4}$. Thus by **(V)** we get $\lambda_1(G_3) \leq \sqrt{3n^{3/4}} = o(\sqrt{\Delta})$.

Let u be a vertex of $V_3 \cup V_4$ which has at least $2n^{2/5}$ neighbors in the set V_4 . Since every vertex in V_4 has at least one neighbor in V_1 we obtain that there at least $2n^{2/5}$ simple paths of length two from u to V_1 . On the other hand we know that there are at most two such paths between any pair of distinct vertices. This implies that u has at least $n^{2/5}$ vertices within distance two whose degree is at least $n^{2/3}$. Using Lemma 2.2 (i) with $a = 2/5$ and $b = 2/3$ we conclude that almost surely there is no such vertex u . Now let G_4 be a subgraph induced by the set V_4 and let G_5 be the bipartite graph with parts (V_3, V_4) . By the above discussion, the maximum degree of G_4 is at most $2n^{2/5}$, implying $\lambda_1(G_4) \leq 2n^{2/5} = o(\sqrt{\Delta})$. We also know that every vertex from V_3 has at most $2n^{2/5}$ neighbors in V_4 and every vertex in V_4 has at most $np + np/\log n \leq n^{5/9}$ neighbors in V_3 . Therefore by **(IV)** we obtain that $\lambda_1(G_5) \leq \sqrt{2n^{2/5}n^{5/9}} = \sqrt{2}n^{43/90} = o(\sqrt{\Delta})$.

Finally consider the bipartite subgraph G_6 whose parts are V_2 and $V_3 \cup V_4$. Let X be the set of vertices from $V_3 \cup V_4$ with at least $2n^{2/7}$ neighbors in V_2 and let $Y = V_3 \cup V_4 - X$. Note that $G_6 = G'_6 \cup G''_6$ where G'_6 is a bipartite graph with parts (V_2, X) and G''_6 ,

is a bipartite graph with parts (V_2, Y) . The upper bound on $\lambda_1(G''_6)$ follows immediately from the facts that G''_6 is bipartite, the degree of vertices in V_2 is bounded by $n^{2/3}$ and, by definition, every vertex in Y has at most $2n^{2/7}$ neighbors in V_2 . Therefore $\lambda_1(G''_6) \leq \sqrt{2n^{2/7}n^{2/3}} = \sqrt{2}n^{10/21} = o(\sqrt{\Delta})$. To bound $\lambda_1(G'_6)$, note that almost surely every vertex in V_2 has at most $n^{3/7}$ neighbors in X . Indeed, let $u \in V_2$ be the vertex with more than $n^{3/7}$ neighbors in X . Since every neighbor of u in X has at least $2n^{2/7} - 1$ additional neighbors in V_2 different from u we obtain that there are at least $(2n^{2/7} - 1)n^{3/7} = (2 + o(1))n^{5/7}$ simple paths of length two from u to V_2 . On the other hand we know that there are at most two such paths between any pair of distinct vertices. This implies that u has at least $(1 + o(1))n^{5/7}$ vertices of V_2 within distance two. By definition, every vertex of V_2 has at least $np + np/\log n$ neighbors in G . Therefore using Lemma 2.2 (ii) with $a = 1/22$ we conclude that almost surely there is no such vertex u . Now the upper bound on $\lambda_1(G'_6)$ can be obtained using that G'_6 is bipartite, the degree of vertices in X is bounded by $np + np/\log n \leq n^{5/9}$ and that every vertex in V_2 has at most $2n^{3/7}$ neighbors in X . Indeed, by **(IV)** $\lambda_1(G'_6) \leq \sqrt{2n^{3/7}n^{5/9}} = \sqrt{2}n^{31/63} = o(\sqrt{\Delta})$, and hence $\lambda_1(G_6) \leq \lambda_1(G'_6) + \lambda_1(G''_6) = o(\sqrt{\Delta})$.

From the above definitions it is easy to check that $G = \cup_{i=1}^6 G_i$. Hence using our estimates on the largest eigenvalues of graphs G_i we obtain the desired upper bound on $\lambda_1(G)$, as follows:

$$\begin{aligned} \lambda_1(G) &\leq \lambda_1(G_1 \cup G_2) + \sum_{i \geq 3} \lambda_1(G_i) \leq (1 + o(1)) \max\left(\sqrt{\Delta(G)}, np\right) + o(\sqrt{\Delta(G)}) \\ &= (1 + o(1)) \max\left(\sqrt{\Delta(G)}, np\right). \end{aligned}$$

This completes the proof of the third case. Now to finish the proof of the theorem it remains to deal with the last very simple case when the random graph is very sparse.

Case 4. Let $p \leq e^{-\log^4 n}$. For every integer $k \geq 1$ denote by Y_k the number of connected components with k edges. It is not difficult to see that $\mathbf{E}Y_k \leq O(2^n k! n^k p^k)$. Indeed, we can pick the first vertex in the connected component in 2^n ways. Suppose we already know the first $1 \leq t \leq k$ vertices of the component. Then these vertices are incident to at most tn edges of the n -cube and therefore we can pick the next edge only in at most tn ways. This gives at most $2^n \prod_{t=1}^k tn = 2^n k! n^k$ ways to pick the edges of the connected component.

First consider the case when p is not exponentially small. Then, by Lemma 2.1 we have that almost surely $\Delta(G) = (1 + o(1))\kappa(n)$, where $\kappa(n) = \max\{k : 2^n \binom{n}{k} p^k (1 - p)^{n-k} \geq 1\}$ and $\kappa(n)$ tends to infinity together with n . Let $k_0 = \kappa(n) + \kappa(n)/\log \kappa(n)$. Then it is easy to check that $\mathbf{E}Y_{k_0} = o(1)$ and therefore, by Markov's inequality, almost surely $G(Q^n, p)$ contains no connected component with more than k_0 edges. Since the largest eigenvalue of G is the maximum of the eigenvalues of its connected components and the largest eigenvalue of a component with k_0 edges is not greater than $\sqrt{k_0}$ (see, parts **(II)** and **(III)** of Proposition 3.1), we obtain that

$$\lambda_1(G) \leq \sqrt{k_0} \leq \sqrt{\kappa(n)} + \sqrt{\kappa(n)/\log \kappa(n)} = (1 + o(1))\sqrt{\Delta(G)}.$$

Next, let $p \leq 2^{-\alpha n}$ for some fixed $\alpha > 0$. If p is not proportional to $2^{-n/k} n^{-1}$, $k = 1, 2, 3, \dots$, then it follows from part (iii) of Lemma 2.1 that with probability going to

one the maximum degree of $G(Q^n, p)$ is $\kappa(n) = \left\lceil \frac{n \log 2}{\log(p^{-1}) - \log n} \right\rceil$. Note that in this case $\kappa(n)$ is a constant and it is easy to check that $\mathbf{E}Y_{\kappa(n)+1} \leq O(2^n n^{\kappa(n)+1} p^{\kappa(n)+1}) = o(1)$. Thus, by Markov's inequality, there are no connected components with more than $\kappa(n)$ edges. Since the largest eigenvalue of G is the maximum of the eigenvalues of its connected components and the largest eigenvalue of a component with k edges is not greater than \sqrt{k} (and is equal to \sqrt{k} only if the component is a star on $k+1$ vertices), we obtain that a.s. $\lambda_1(G) = \sqrt{\kappa(n)} = \sqrt{\Delta(G)}$.

Finally if $p(n)$ is proportional to $2^{-n/k} n^{-1}$, $k = 1, 2, 3, \dots$, then by part (ii) of Lemma 2.1 almost surely $\Delta(G) \in \{k-1, k\}$ and again one can check that $\mathbf{E}Y_{k+1}$ is exponentially small. Using Markov's inequality, as before, we conclude that there are no connected components with more than k edges. Therefore a.s. $\lambda_1(G)$ is either $\sqrt{\Delta(G)}$ or $\sqrt{\Delta(G)+1}$. This completes the proof of the theorem.

4. Concluding Remarks

There are several other important questions that are beyond the reach of the presented technique. The most fundamental is perhaps the local statistics of the eigenvalues, in particular the local statistics near the edge of the spectrum. For results in this direction for other random matrix models we refer the reader to [17, 18, 16]. A recent result of Alon, Krivelevich and Vu [3] states that the deviation of the first, second, etc. largest eigenvalues from its mean is at most of order of $O(1)$. Unfortunately our results give only the leading term of the mean.

A second, perhaps even more difficult question is whether the local behavior of the eigenvalues is sensitive to the details of the distribution of the matrix entries of A . We refer the reader to [15, 6, 16, 10] for the results of that nature for unitary invariant and Wigner random matrices.

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